

Recitation 9

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Problem 1: [YG] Problem 9.1.3 a-c

Solution:

(a) For $0 \leq x \leq 1$,

$$f_X(x) = \int_{-\infty}^{\infty} f_{X,Y}(x,y) dy = \int_x^1 2 dy = 2(1-x). \quad (1)$$

The complete expression of the PDF of X is

$$f_X(x) = \begin{cases} 2(1-x) & 0 \leq x \leq 1, \\ 0 & \text{otherwise.} \end{cases} \quad (2)$$

(b) The blind estimate of X is

$$\hat{X}_B = E[X] = \int_0^1 2x(1-x) dx = \left(x^2 - \frac{2x^3}{3}\right) \Big|_0^1 = \frac{1}{3}. \quad (3)$$

(c) First we calculate

$$P[X > 1/2] = \int_{1/2}^1 f_X(x) dx = \int_{1/2}^1 2(1-x) dx = (2x - x^2) \Big|_{1/2}^1 = \frac{1}{4}. \quad (4)$$

Now we calculate the conditional PDF of X given $X > 1/2$.

$$f_{X|X>1/2}(x) = \begin{cases} \frac{f_X(x)}{P[X>1/2]} & x > 1/2, \\ 0 & \text{otherwise,} \end{cases} = \begin{cases} 8(1-x) & 1/2 < x \leq 1, \\ 0 & \text{otherwise.} \end{cases} \quad (5)$$

The minimum mean square error estimate of X given $X > 1/2$ is

$$E[X|X > 1/2] = \int_{-\infty}^{\infty} x f_{X|X>1/2}(x) dx \quad (6)$$

$$= \int_{1/2}^1 8x(1-x) dx = \left(4x^2 - \frac{8x^3}{3}\right) \Big|_{1/2}^1 = \frac{2}{3}. \quad (7)$$

Problem 2: [YG] Problem 6.6.1

Solution:

We know that the waiting time, W is uniformly distributed on $[0,10]$ and therefore has the following PDF.

$$f_W(w) = \begin{cases} 1/10 & 0 \leq w \leq 10 \\ 0 & \text{otherwise} \end{cases} \quad (1)$$

We also know that the total time is 3 milliseconds plus the waiting time, that is $X = W + 3$.

- (a) The expected value of X is $E[X] = E[W + 3] = E[W] + 3 = 5 + 3 = 8$.
- (b) The variance of X is $\text{Var}[X] = \text{Var}[W + 3] = \text{Var}[W] = 25/3$.
- (c) The expected value of A is $E[A] = 12E[X] = 96$.
- (d) The standard deviation of A is $\sigma_A = \sqrt{\text{Var}[A]} = \sqrt{12(25/3)} = 10$.
- (e) $P[A > 116] = 1 - \Phi\left(\frac{116-96}{10}\right) = 1 - \Phi(2) = 0.02275$.
- (f) $P[A < 86] = \Phi\left(\frac{86-96}{10}\right) = \Phi(-1) = 1 - \Phi(1) = 0.1587$

Problem 2

- a X_i takes the value 1 with probability p and the value 0 with probability $1-p$. Let $S_n = \sum_{i=1}^n X_i$ where the X_i 's are mutually independent and identically distributed. Compute (i) the probability mass-function (pmf), (ii) the mean, and (iii) the variance of S_n .
- b Suppose X_i 's are binary digits with probabilities $p(1) = 0.005$ and $p(0) = 0.995$. The digits are taken 100 at a time and a codeword is provided for every sequence of 100 digits containing three or fewer 1's. (i) Calculate the probability of observing a sequence for which no codeword has been assigned. (ii) Use Chebyshev's inequality to bound the probability of observing a sequence for which no codeword has been assigned. Compare the bound with the actual probability computed in (i).

Solution:

- a (i) The sum of independent Bernoulli random variables is a Binomial random variable.

$$P(S_n = k) = \binom{n}{k} p^k (1-p)^{n-k}$$

(ii)

$$E[S_n] = E\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n E[X_i] = nE[X_1] = np$$

(iii)

$$\text{Var}[S_n] = \text{Var}\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n \text{Var}[X_i] = n\text{Var}[X_1] = np(1-p)$$

b (i) The probability that a 100-bit sequence has three or fewer ones is

$$\sum_{i=0}^3 \binom{100}{i} (0.005)^i (0.995)^{100-i} = 0.99833$$

Thus the probability that the sequence that is generated cannot be encoded is $1 - 0.99833 = 0.00167$.

(ii) In the case of a random variable S_n that is the sum of n i.i.d random variables X_1, \dots, X_n , Chebyshev's inequality states that

$$P(|S_n - n\mu| \geq a) \leq \frac{n\sigma^2}{a^2}$$

where μ and σ^2 are the mean and variance of the X 's (therefore $n\mu$ and $n\sigma^2$ are the mean and variance of S_n). In this problem, $n = 100$, $\mu = p = 0.005$ and $\sigma^2 = p(1-p) = (0.005)(0.995)$. Further note that $S_{100} \geq 4$ if and only if $|S_{100} - 100(0.005)| \geq 3.5$, so we should choose $a = 3.5$. Then,

$$P(S_{100} \geq 4) \leq \frac{100(0.005)(0.995)}{(3.5)^2} \approx 0.04061$$

This bound is about 24 times bigger than the actual probability 0.00167. This shows that the Chebyshev inequality does not truly reflect the rate of decay of probability with blocklength n .