

SE524/EC524 Optimization Theory and Methods

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Lecture 3: Outline

- ➊ Corner Points:
 - ➊ Extreme point.
 - ➋ Vertex.
 - ➌ Basic Feasible Solution (BFS).
- ➋ Theorem: Extreme point \Leftrightarrow Vertex \Leftrightarrow BFS.
- ➌ BFS in standard form polyhedra.
- ➍ Degeneracy is standard form polyhedra.

Extreme Points

Definition

Let P a polyhedron. A vector $\mathbf{x} \in P$ is an **extreme point** if $\nexists \mathbf{y}, \mathbf{z} \in P$ and $\lambda \in (0, 1)$ s.t. \mathbf{x} can be written as a convex combination of \mathbf{y}, \mathbf{z} .

Vertex

Definition

A vector $\mathbf{x} \in P$ is a **vertex** if $\exists \mathbf{c}$ s.t. \mathbf{x} is the unique minimizer of $\mathbf{c}'\mathbf{y}$ over P .

Basic Feasible Solution (BFS)

We call **active** or **binding** the constraints that are satisfied with equality.

Definition

Let $\mathbf{x}^* \in \mathbb{R}^n$, and a polyhedron P .

- 1 \mathbf{x}^* is a **basic solution** if
 - 1 All equality constraints are active.
 - 2 Let $I = \{i \mid \mathbf{a}_i^T \mathbf{x}^* = \mathbf{b}_i\}$. The subspace spanned by $\{\mathbf{a}_i \mid i \in I\}$ is all of \mathbb{R}^n .
- 2 If \mathbf{x}^* is feasible then is called a **basic feasible solution (bfs)**.

A basic solution \mathbf{x}^* is called **degenerate** if $|I| > n$, and **nondegenerate** otherwise.

Equivalence of corner point definitions

Theorem

Extreme point \Leftrightarrow Vertex \Leftrightarrow BFS.

Proof:

- 1 Vertex \Rightarrow Extreme point.
- 2 Extreme point \Rightarrow BFS.
- 3 BFS \Rightarrow Vertex.

□

- Two basic solutions are **adjacent** if there are $n - 1$ linearly independent constraints active at both of them.
- If two bfs are adjacent, the line segment that connects them is called an **edge**.

BFS in standard form polyhedra

Let $P = \{x \mid Ax = b, x \geq 0\}$ nonempty.

WLOG assume A has m linearly independent rows.

Example

$$\begin{aligned}x_1 + 2x_2 &= 2 \\2x_1 + x_2 &= 4 \\2x_1 + 4x_2 &= 4\end{aligned}$$

For a BFS we need n active "lin. ind." constraints: $Ax = b$ provides m . Hence, $n - m$ of $x \geq 0$ should be active.

BFS in standard form polyhedra: Construction

Pick m linearly independent columns of A (wlog $1, 2, \dots, m$).

$$m \left\{ \begin{array}{c} \overbrace{B}^m \\ \left| \right. \\ \overbrace{N}^{n-m} \end{array} \right\} \begin{bmatrix} x_B \\ x_N \end{bmatrix} = b \quad \Rightarrow \quad \begin{cases} x_B = B^{-1}b \\ x_N = 0 \end{cases}$$

- $x^* = (x_B, x_N)$ is a **basic solution**.
- If $x^* \geq 0$ it is a **basic feasible solution**.
- The elements of x_B are called **basic variables**.
- The elements of x_N are called **nonbasic variables**.
- B is called a **basis matrix**, its columns form a **basis** (of \mathbb{R}^m).
- Two bases that share $m - 1$ columns are called **adjacent** and they correspond to adjacent bfs.
- A basis **uniquely** determines a bfs, **but** different bases may lead to the same bfs.

Degeneracy is standard form polyhedra

Definition

When the set of active constraints at a BFS has more than n elements then the BFS is **degenerate**.

In **standard form**: $\mathbf{Ax} = \mathbf{b}$ are m active constraints. If more than $n - m$ of $\mathbf{x} \geq \mathbf{0}$ active then **degenerate bfs**.

The following statements are equivalent and imply **degeneracy**:

- More than $n - m$ variables = 0.
- Some basic variables = 0.